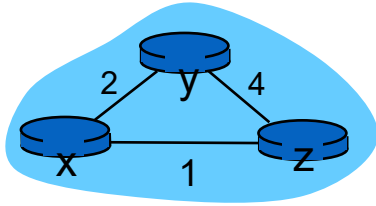


Advanced Computer Networks midterm (110/4)

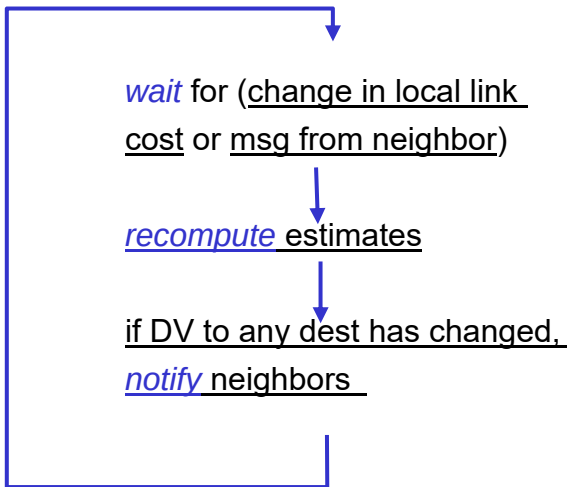
只寫答案而沒有解釋說明，扣一半分數

1. (a) Describe the flow of distance vector routing algorithm. (6%)
(b) List three tables of node X, Y and Z with the distance vector algorithm, from the time when three nodes are initialized to the time three tables are stabilized. (table 一行 3% (x, y, z 各看自己那列 1%), 共 6%。數值有變動時，要寫出公式 2%，共 4% => 10%, 16% total)



Ans:

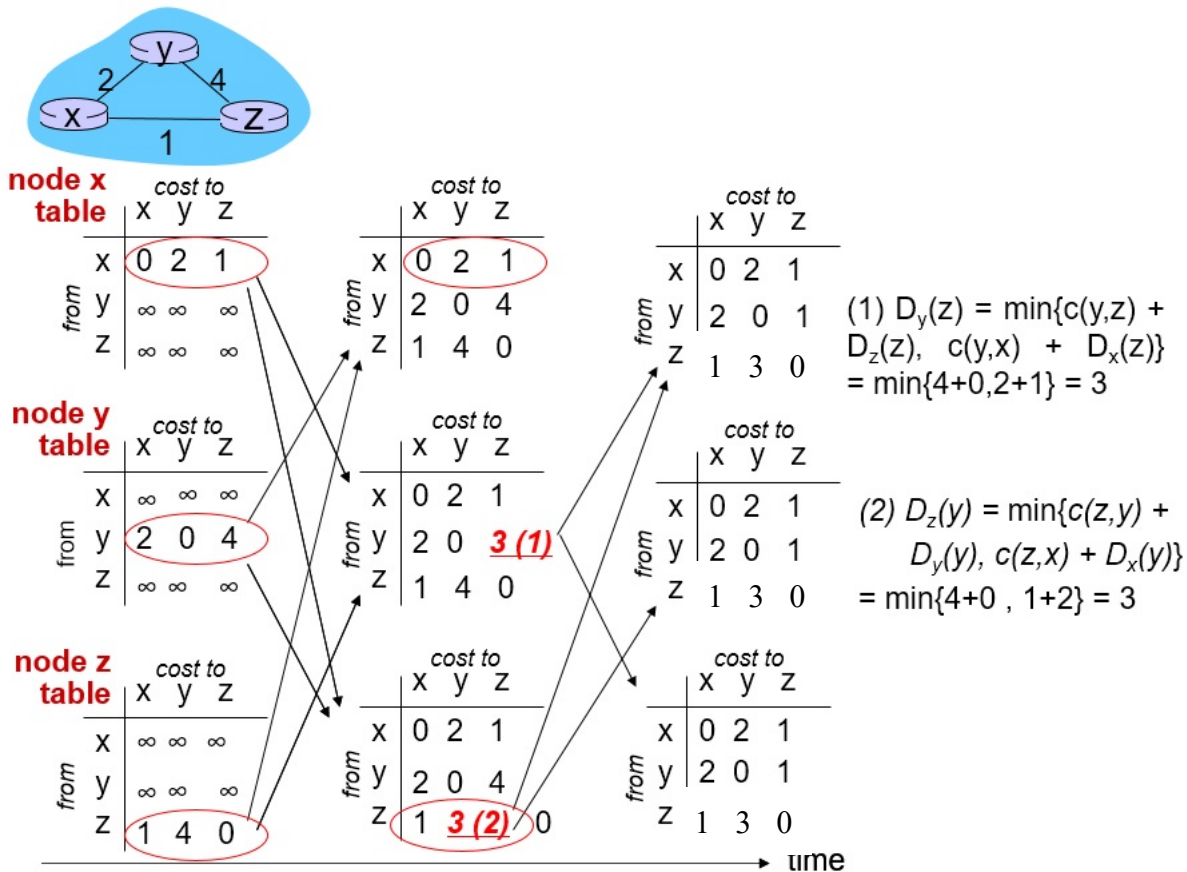
(a)



(2% per item)

(b)

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2. What kind of routing algorithm does OSPF use? List its three steps. (8%)

Ans:

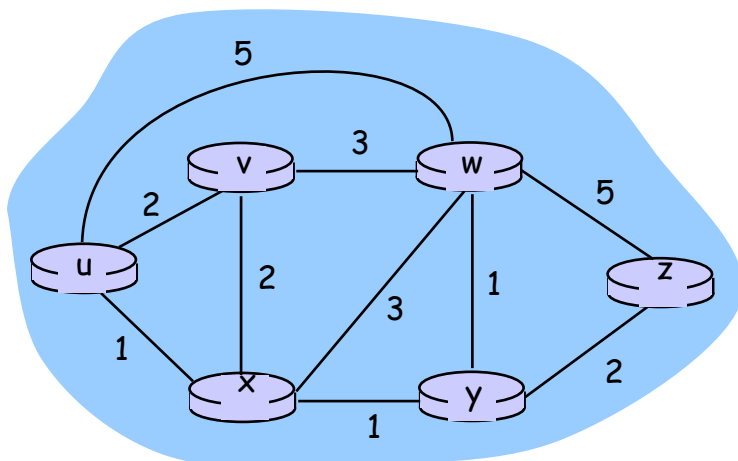
uses Link State algorithm (2%)

LS packet dissemination (2%)

topology map at each node (2%)

route computation using Dijkstra's algorithm (2%)

3. Use Dijkstra's shortest-path algorithm to compute the shortest path from the source node to all other network nodes. (a) Show how the algorithm works by computing a table. (公式錯一個扣 1%，含過程的表格共 10%，扣完為止) (b) Show the forwarding table of the source node. (5%) (note: the source node 選法：學號最後一位除以 5 的餘數：0=>z, 1=>y, 2=>x, 3=>w, 4=>v, cost 數值相同時，優先選字母順序較前者；數值有變動時，要寫出公式) (15% total)



Ans: (除 step 之外，一欄 1 分)

以 v 為起點

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Step	N'	D(u), p(u)	D(w), p(w)	D(x), p(x)	D(y), p(y)	D(z), p(z)
0	v	<u>2, v</u>	3, v	2, v	∞	∞
1	vu		3, v	<u>2, v</u>	∞	∞
2	vux		<u>3, v</u>		3, x #1	∞
3	vuxw				<u>3, x</u>	8, w #2
4	vuxwy					<u>5, y</u> #3
5	vuxwyz					

#1. $D(y) = \min(D(y), D(x) + c(x, y)) = \min(\infty, 2 + 1) = 3$ (1%)

#2. $D(z) = \min(D(z), D(w) + c(w, z)) = \min(\infty, 3 + 5) = 8$

#3. $D(z) = \min(D(z), D(y) + c(y, z)) = \min(\infty, 3 + 2) = 5$

Forwarding table (一列 1 分)

Destination	Next hop (output link)
u	u
w	w
x	x
y	x
z	x

以 w 為起點

Step	N'	D(u), p(u)	D(v), p(v)	D(x), p(x)	D(y), p(y)	D(z), p(z)
0	w	5, w	3, w	3, w	<u>1, w</u>	5, w
1	wy	5, w	3, w	<u>2, y</u> #1		3, y #2
2	wyx	<u>3, x</u> #3	3, w			3, y
3	wyxu		<u>3, w</u>			3, y
4	wyxuv					<u>3, y</u>
5	wyxuvz					

#1. $D(x) = \min(D(x), D(y) + c(y, x)) = \min(3, 1 + 1) = 2$

#2. $D(z) = \min(D(z), D(y) + c(y, z)) = \min(5, 1 + 2) = 3$

#3. $D(u) = \min(D(u), D(x) + c(x, u)) = \min(5, 2 + 1) = 3$

Forwarding table

Destination	Next hop (output link)
u	y
v	v
x	y
y	y
z	y

以 x 為起點

Step	N'	D(u), p(u)	D(v), p(v)	D(w), p(w)	D(y), p(y)	D(z), p(z)
0	x	<u>1, x</u>	2, x	3, x	1, x	∞

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1	xu		2, x	3, x	<u>1, x</u>	∞
2	xuy		<u>2, x</u>	2, y #1		3, y #2
3	xuyv			<u>2, y</u>		3, y
4	xuyvw					<u>3, y</u>
5	xuyvwz					

#1. $D(w) = \min(D(w), D(y) + c(y, w)) = \min(3, 1 + 1) = 2$

#2. $D(z) = \min(D(z), D(y) + c(y, z)) = \min(\infty, 1 + 2) = 3$

Forwarding table

Destination	Next hop (output link)
u	u
v	v
w	y
y	y
z	y

以 y 為起點

Step	N'	D(u), p(u)	D(v), p(v)	D(w), p(w)	D(x), p(x)	D(z), p(z)
0	y	∞	∞	<u>1, y</u>	1, y	2, y
1	yw	6, w #1	4, w #2		<u>1, y</u>	2, y
2	ywx	<u>2, x</u> #3	3, x #4			2, y
3	ywxu		3, x			<u>2, y</u>
4	ywxuz		<u>3, x</u>			
5	ywxuzv					

#1. $D(u) = \min(D(u), D(w) + c(w, u)) = \min(\infty, 1 + 5) = 6$

#2. $D(v) = \min(D(v), D(w) + c(w, v)) = \min(\infty, 1 + 3) = 4$

#3. $D(u) = \min(D(u), D(x) + c(x, u)) = \min(\infty, 1 + 1) = 2$

#4. $D(v) = \min(D(v), D(x) + c(x, v)) = \min(\infty, 1 + 2) = 3$

Forwarding table

Destination	Next hop (output link)
u	x
v	x
w	w
x	x
z	z

以 z 為起點

Step	N'	D(u), p(u)	D(v), p(v)	D(w), p(w)	D(x), p(x)	D(y), p(y)
0	z	∞	∞	5, z	∞	<u>2, z</u>
1	zy	∞	∞	<u>3, y</u> #1	3, y #2	

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2	zyw	8, w #3	6, w #4		<u>3, y</u>	
3	zywx	<u>4, x</u> #5	5, x #6			
4	zywxu		<u>5, x</u>			
5	zywxuv					

#1. $D(w) = \min(D(w), D(y) + c(y, w)) = \min(5, 2+1) = 3$

#2. $D(x) = \min(D(x), D(y) + c(y, x)) = \min(\infty, 2+1) = 3$

#3. $D(u) = \min(D(u), D(w) + c(w, u)) = \min(\infty, 3+5) = 8$

#4. $D(v) = \min(D(v), D(w) + c(w, v)) = \min(\infty, 3+3) = 6$

#5. $D(u) = \min(D(u), D(x) + c(x, u)) = \min(\infty, 3+1) = 4$

#6. $D(v) = \min(D(v), D(x) + c(x, v)) = \min(\infty, 3+2) = 5$

Forwarding table

Destination	Next hop (output link)
u	y
v	y
w	y
x	y
y	y

4. 針對 192.1.2.1 這個 IP address，(以十進位表示，要寫完整)

- 這一個 IP 屬於哪個 Class 的網路？以二進位說明(1%)其所屬的 IP 網路表示為何？(1%) 可用 IP 範圍？(2%) 共有幾個 IP 可用？(1%) mask 的值為何？(1%)
- 將此 IP 網路分成 9 個 subnet，subnet mask 的值為何？(1%) 請列出第 6 個 subnet 的網路表示法為何？(1%) 可用 IP 範圍？(2%) 共有幾個 IP 可用？(1%) (11% total)

Ans:

a.

192.1.2.1 的二進位表示法為 11000000.00000001.00000010.xxxxxxxx，由前三個 bits 110 可判斷為 **Class C** 的 IP (2%)

此 IP 所屬於的 Class C 的網路表示為 192.1.2.0 (2%)

所有 host ID 部分的 8 個 bit 的 x 不可以全為 0 或 1，因此第一個可用 host ID 為

11000000.00000001.00000010.00000001 = 192.1.2.1 (2%)

最後一個可用 host ID 為 11000000.00000001.00000010.11111110 = 192.1.2.254 (2%)，共有 $2^8 - 2 = 254$ 個可用 host ID (2%)

Mask: 255.255.255.0 (2%)

b.

將此 Class C 網路分成 9 個 subnet，最少需要 $9 \leq 2^4$ ，subnet mask 的值 → 需要 host ID 的前 4 個 bits 當作 subnet ID。所以新的 subnet mask 是由原本 Class C 的 default subnet mask 255.255.255.0 來改，改成 255.255.255.11110000 = 255.255.255.240 (2%)

subnet 的 ID 要從此 Class C Network ID 11000000.00000001.00000010.xxxxxxxx 來改，需要 ID 的前 4 個 bits 當作 subnet ID。因此第 6 個 subnet ID (從 000 開始，第六個是 0101) 為

11000000.00000001.00000010.01010000 → 192.1.2.80 (2%)，因此第一個可用 host ID 為

11000000.00000001.00000010.01010001 = 192.1.2.81 (2%)

最後一個可用 host ID 為 11000000.00000001.00000010.01011110 = 192.1.2.94 (2%)，共有 $2^4 - 2 = 14$ 個可用 host ID (2%)

- (a) What is the goal of DHCP? (2%) (b) List four steps of DHCP (4%) (c) 使用 DHCP 自動設定 IP 時，會取得哪四項資訊，電腦可以上網？(4%) (10% total)

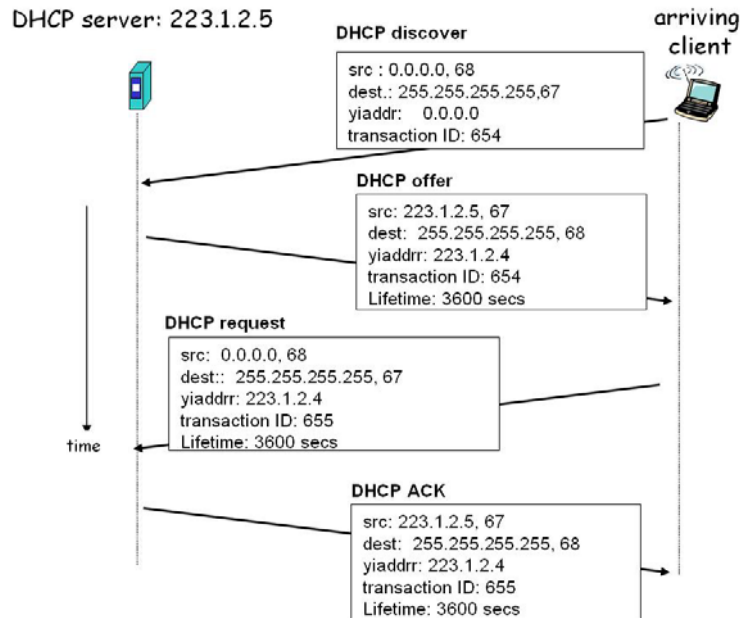
Advanced Computer Networks midterm (110/4)

Ans:

(a) Goal: allow host to *dynamically* obtain its IP address from network server when it joins network (2%)

(b) Flow: (4%)

- host broadcasts “DHCP discover” msg (1%)
- DHCP server responds with “DHCP offer” msg (1%)
- host requests IP address: “DHCP request” msg (1%)
- DHCP server sends address: “DHCP ack” msg (1%)



(c) IP address, subnet mask, default gateway, DNS server (4%)

6. Describe how Ethernet uses CSMA/CD with exponential backoff (要寫出碰撞後如何動作) in detail (12%)

Ans:

- ▶ adapter doesn't transmit if it senses that some other adapter is transmitting, that is, **carrier sense** (2%)
- ▶ transmitting adapter aborts when it senses that another adapter is transmitting, that is, **collision detection** (2%)
- ▶ Before attempting a retransmission, adapter waits a random time, that is, **random access with Exponential Backoff**. (2%)
 - ▶ first collision: choose K from {0,1}; delay is K · 512 bit transmission times (2%)
 - after second collision: choose K from {0,1,2,3}... (2%)
 - after m collisions, choose K from {0,1,2,3,4,...,2^m-1} (2%)

7. (a) List three types of multiple access protocols and describe how they work briefly. (9%) (b) Classify FDMA, Token Passing and CSMA/CD into one of the type to whom they belong. (6%) (15% total)

Ans:

(a) Three broad classes:

a. Channel Partitioning (2%)

- ▶ divide channel into smaller “pieces” (time slots, frequency, code) to node for exclusive use (1%)

b. Random Access (2%)

- ▶ channel not divided, allow collisions (1%)

c. “Taking turns” (2%)

- ▶ Nodes take turns, but nodes with more to send can take longer turns (1%)

(b) Channel Partitioning: FDMA; (2%)

Random Access: CSMA/CD (2%)

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“Taking turns”: Token Passing (2%)

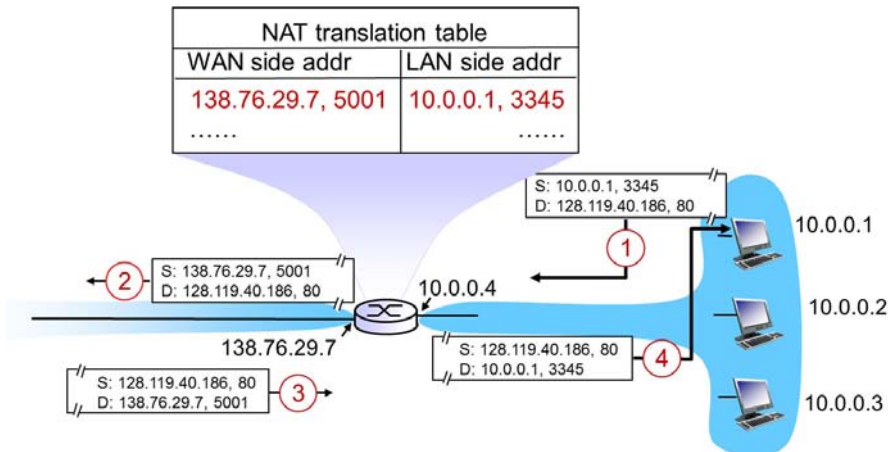
8. Consider the CRC generator, $G=1001$, and suppose that D has the value 10101010000 . What is the value of R ? (要寫出運算過程 6%, 8% total)

Ans:

If we divide 1001 into 10101010000 , we get 10111101 (過程 6%), with a remainder of $R = 101$ (2%).

$$\begin{array}{r}
 10111101 \\
 1001 \overline{) 10101010000} \\
 \underline{1001} \\
 1110 \\
 \underline{1001} \\
 1111 \\
 \underline{1001} \\
 1100 \\
 \underline{1001} \\
 1010 \\
 \underline{1001} \\
 1100 \\
 \underline{1001} \\
 101
 \end{array}$$

9. Explain four NAT operations with this figure (用圖上的數值說明 5%)



Ans:

NAT router must: (8%)

- outgoing datagrams: replace (source IP address, port #)=(10.0.0.1, 3345) (1%) of every outgoing datagram to (NAT IP address, new port #)=(138.76.29.7, 5001) (1%)
- remember (in NAT translation table) every (source IP address, port #)=(10.0.0.1, 3345) to (NAT IP address, new port #)=(138.76.29.7, 5001) translation pair (1%)
- incoming datagrams: replace (NAT IP address, new port #)=(138.76.29.7, 5001) (1%) in dest fields of every incoming datagram with corresponding (source IP address, port #)=(10.0.0.1, 3345) (1%) stored in NAT table