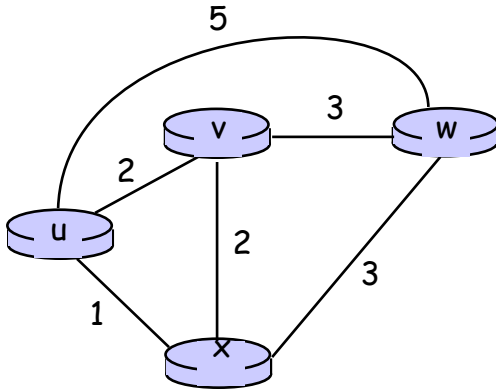


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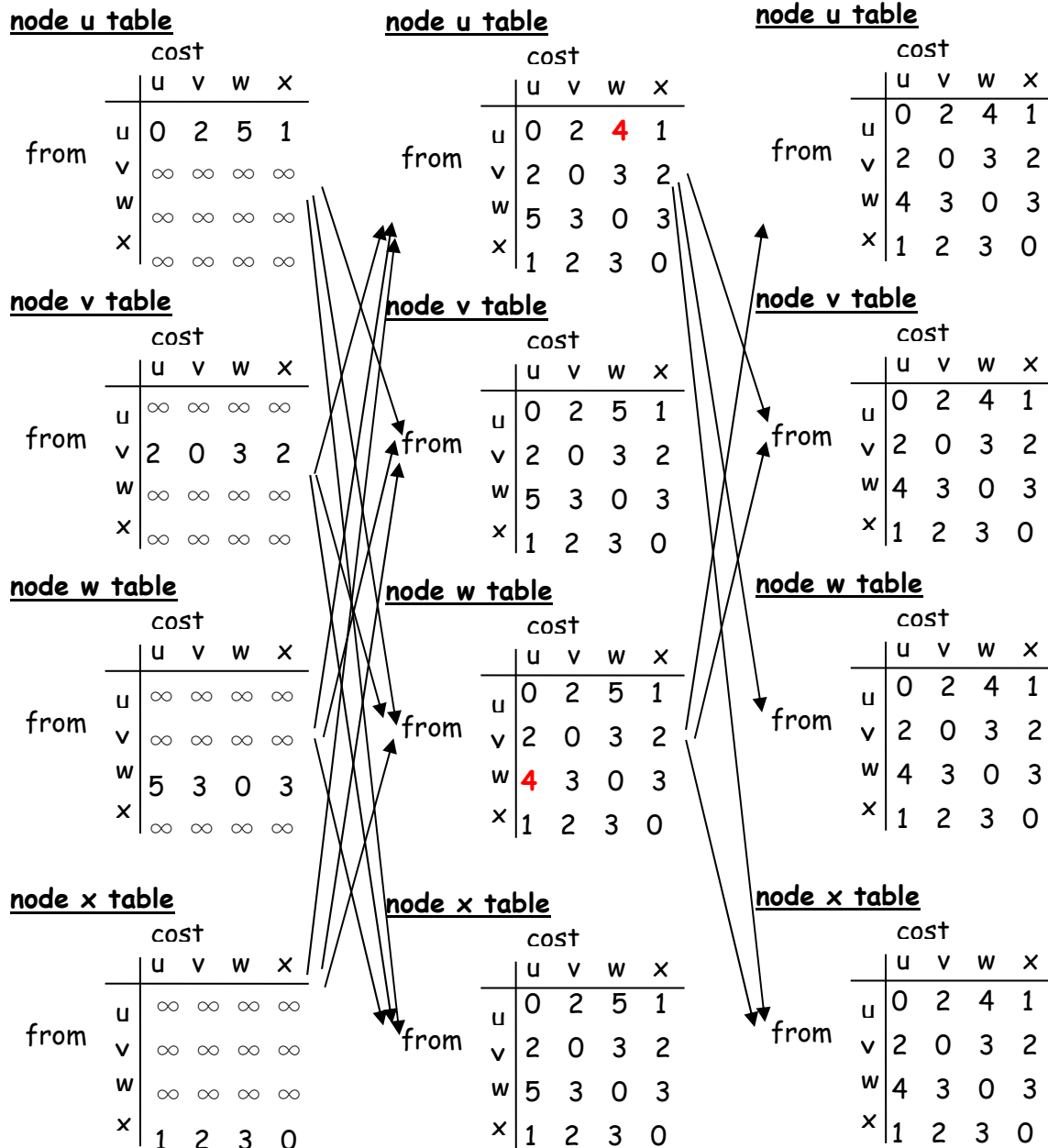
只寫答案而沒有解釋說明，扣一半分數

1. List four tables of node u, v, w and x with the distance vector algorithm, from the time when four nodes are initialized to the time four tables are stabilized. (table 一列 2% , u, v, w and x 各看自己那列 8% 。寫出從初始化步驟到第二步驟時每個 node 到其他 node 之 distance vector 計算的三條公式，如 u 為 $D_u(v)$, $D_u(w)$, $D_u(x)$ 6% ，畫出傳送 distance vector 的箭頭 2% ， 16% total)



Ans:

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U:

$$D_u(v) = \min \{c(u,x) + D_x(v), c(u,w) + D_w(v), c(u,v) + D_v(v)\}$$

$$= \min \{1+2, 5+3, 2+0\} = 2$$

$$D_u(w) = \min \{c(u,x) + D_x(w), c(u,w) + D_w(w), c(u,v) + D_v(w)\}$$

$$= \min \{1+3, 5+0, 2+3\} = 4$$

$$D_u(x) = \min \{c(u,x) + D_x(x), c(u,w) + D_w(x), c(u,v) + D_v(x)\}$$

$$= \min \{1+0, 5+3, 2+2\} = 1$$

u's Forwarding table

Destination	Next hop (output link)
v	v
w	x
x	x

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V:

$$D_v(u) = \min\{c(v,u) + D_u(u), c(v,w) + D_w(u), c(v,x) + D_x(u)\} \\ = \min\{2+0, 3+5, 2+1\} = 2$$

$$D_v(w) = \min\{c(v,x) + D_x(w), c(v,w) + D_w(w), c(v,u) + D_u(w)\} \\ = \min\{2+3, 3+0, 2+5\} = 3$$

$$D_v(x) = \min\{c(v,x) + D_x(x), c(v,w) + D_w(x), c(v,u) + D_u(x)\} \\ = \min\{2+0, 3+3, 2+1\} = 2$$

v's Forwarding table

Destination	Next hop (output link)
u	u
w	w
x	x

W:

$$D_w(u) = \min\{c(w,u) + D_u(u), c(w,v) + D_v(u), c(w,x) + D_x(u)\} \\ = \min\{5+0, 3+2, 3+1\} = 4$$

$$D_w(v) = \min\{c(w,v) + D_v(v), c(w,x) + D_x(v), c(w,u) + D_u(v)\} \\ = \min\{3+0, 3+2, 5+2\} = 3$$

$$D_w(x) = \min\{c(w,x) + D_x(x), c(w,v) + D_v(x), c(w,u) + D_u(x)\} \\ = \min\{3+0, 3+2, 5+1\} = 3$$

w's Forwarding table

Destination	Next hop (output link)
u	x
v	v
x	x

X:

$$D_x(u) = \min\{c(x,u) + D_u(u), c(x,v) + D_v(u), c(x,w) + D_w(u)\} \\ = \min\{1+0, 2+2, 3+5\} = 1$$

$$D_x(v) = \min\{c(x,v) + D_v(v), c(x,w) + D_w(v), c(x,u) + D_u(v)\} \\ = \min\{2+0, 3+3, 1+2\} = 2$$

$$D_x(w) = \min\{c(x,w) + D_w(w), c(x,v) + D_v(w), c(x,u) + D_u(w)\} \\ = \min\{3+0, 2+3, 1+5\} = 3$$

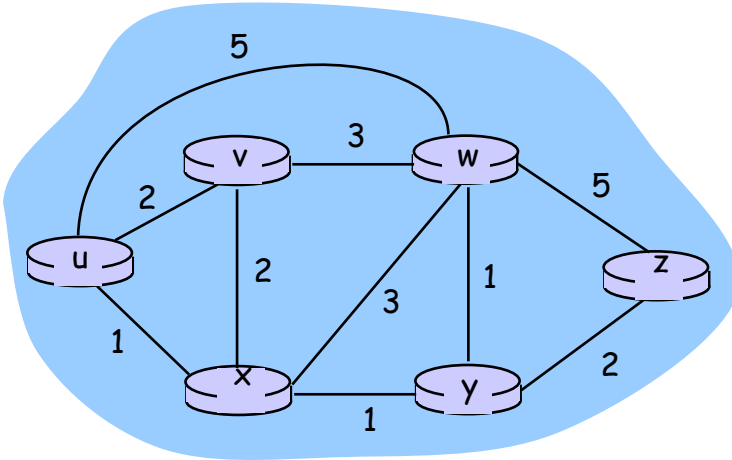
x's Forwarding table

Destination	Next hop (output link)
u	u
v	v
w	w

2. Use Dijkstra's shortest-path algorithm to compute the shortest path from the source node to all other network nodes. (a) Show how the algorithm works by computing a table. (公式錯一個扣 1%，含過程)

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的表格共 10%，扣完為止) (b) Show the forwarding table of the source node. (5%) (note: the source node 選法：學號最後一位除以 5 的餘數：0=>z, 1=>y, 2=>x, 3=>w, 4=>v, cost 數值相同時，優先選字母順序較前者；數值有變動時，要寫出公式) (15% total)



Ans: (除 step 之外，一欄 1 分)

以 v 為起點

Step	N'	D(u), p(u)	D(w), p(w)	D(x), p(x)	D(y), p(y)	D(z), p(z)
0	v	<u>2, v</u>	3, v	2, v	∞	∞
1	vu		3, v	<u>2, v</u>	∞	∞
2	vux		<u>3, v</u>		3, x #1	∞
3	vuxw				<u>3, x</u>	8, w #2
4	vuxwy					<u>5, y</u> #3
5	vuxwyz					

#1. $D(y) = \min(D(y), D(x) + c(x, y)) = \min(\infty, 2 + 1) = 3$ (1%)

#2. $D(z) = \min(D(z), D(w) + c(w, z)) = \min(\infty, 3 + 5) = 8$

#3. $D(z) = \min(D(z), D(y) + c(y, z)) = \min(\infty, 3 + 2) = 5$

Forwarding table (一列 1 分)

Destination	Next hop (output link)
u	u
w	w
x	x
y	x
z	x

以 w 為起點

Step	N'	D(u), p(u)	D(v), p(v)	D(x), p(x)	D(y), p(y)	D(z), p(z)
0	w	5, w	3, w	3, w	<u>1, w</u>	5, w
1	wy	5, w	3, w	<u>2, y</u> #1		3, y #2
2	wyx	<u>3, x</u> #3	3, w			3, y
3	wyxu		<u>3, w</u>			3, y
4	wyxuv					<u>3, y</u>
5	wyxuvz					

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$$\#1.D(x)=\min(D(x),D(y)+c(y,x))=\min(3, 1+1)=2$$

$$\#2.D(z)=\min(D(z),D(y)+c(y,z))=\min(5, 1+2)=3$$

$$\#3.D(u)=\min(D(u),D(x)+c(x,u))=\min(5, 2+1)=3$$

Forwarding table

Destination	Next hop (output link)
u	y
v	v
x	y
y	y
z	y

以 x 為起點

Step	N'	D(u), p(u)	D(v), p(v)	D(w), p(w)	D(y), p(y)	D(z), p(z)
0	x	<u>1, x</u>	2, x	3, x	1, x	∞
1	xu		2, x	3, x	<u>1, x</u>	∞
2	xuy		<u>2, x</u>	2, y #1		3, y #2
3	xuyv			<u>2, y</u>		3, y
4	xuyvw					<u>3, y</u>
5	xuyvwz					

$$\#1.D(w)=\min(D(w),D(y)+c(y,w))=\min(3, 1+1)=2$$

$$\#2.D(z)=\min(D(z),D(y)+c(y,z))=\min(\infty, 1+2)=3$$

Forwarding table

Destination	Next hop (output link)
u	u
v	v
w	y
y	y
z	y

以 y 為起點

Step	N'	D(u), p(u)	D(v), p(v)	D(w), p(w)	D(x), p(x)	D(z), p(z)
0	y	∞	∞	<u>1, y</u>	1, y	2, y
1	yw	6, w #1	4, w #2		<u>1, y</u>	2, y
2	ywx	<u>2, x</u> #3	3, x #4			2, y
3	ywxu		3, x			<u>2, y</u>
4	ywxuz		<u>3, x</u>			
5	ywxuzv					

$$\#1.D(u)=\min(D(u),D(w)+c(w,u))=\min(\infty, 1+5)=6$$

$$\#2.D(v)=\min(D(v),D(w)+c(w,v))=\min(\infty, 1+3)=4$$

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#3. $D(u) = \min(D(u), D(x) + c(x, u)) = \min(\infty, 1 + 1) = 2$

#4. $D(v) = \min(D(v), D(x) + c(x, v)) = \min(\infty, 1 + 2) = 3$

Forwarding table

Destination	Next hop (output link)
u	x
v	x
w	w
x	x
z	z

以 z 為起點

Step	N'	D(u), p(u)	D(v), p(v)	D(w), p(w)	D(x), p(x)	D(y), p(y)
0	z	∞	∞	5, z	∞	<u>2, z</u>
1	zy	∞	∞	<u>3, y</u> #1	3, y #2	
2	zyw	8, w #3	6, w #4		<u>3, y</u>	
3	zywx	<u>4, x</u> #5	5, x #6			
4	zywxu		<u>5, x</u>			
5	zywxuv					

#1. $D(w) = \min(D(w), D(y) + c(y, w)) = \min(5, 2 + 1) = 3$

#2. $D(x) = \min(D(x), D(y) + c(y, x)) = \min(\infty, 2 + 1) = 3$

#3. $D(u) = \min(D(u), D(w) + c(w, u)) = \min(\infty, 3 + 5) = 8$

#4. $D(v) = \min(D(v), D(w) + c(w, v)) = \min(\infty, 3 + 3) = 6$

#5. $D(u) = \min(D(u), D(x) + c(x, u)) = \min(\infty, 3 + 1) = 4$

#6. $D(v) = \min(D(v), D(x) + c(x, v)) = \min(\infty, 3 + 2) = 5$

Forwarding table

Destination	Next hop (output link)
u	y
v	y
w	y
x	y
y	y

3. 針對 163.107.172.1 這個 IP address，(以十進位表示，要寫完整過程) (19%)
- 這一個 IP 屬於那個 Class 的網路？以二進位說明(1%) 其所屬的 IP 網路表示法為何？(2%) 可用 IP 範圍？(2%) 共有幾個 IP 可用？(1%) mask 的值為何？(1%)
 - 將此 IP 網路分成 6 subnets，subnet mask 的值為何？(2%) 請列出第 6 個 subnet 的網路表示法(2%) 可用 IP 範圍？(4%) 共有幾個 IP 可用？(1%)
 - 手動設定電腦的網路時，至少要設定哪三個項目的資訊，才可以上網？(3%)

Ans:

- 163.107.172.1 的二進位表示法為 10100011.01101011.XXXXXXXXXX.XXXXXXXXXX，由前兩個 bits 10 可判斷為 Class B 的 IP。(1%)

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此 IP 所屬於的 Class B 的網路表示法為 163.107.0.0 (2%)

所有 host ID 部分的 16 個 bit 的 X 不可以全為 0 或 1，

因此第一個可用 Host ID 為 10100011.01101011.00000000.00000001 = 163.107.0.1 (1%)

最後一個可用 Host ID 為 10100011.01101011.11111111.11111110 = 163.107.255.254 (1%)

->共有 $2^{16}-2=65534$ 個可用 Host ID (1%)

Mask: 255.255.0.0 (1%)

b.

將此 Class B 網路分成 6 個 subnet，最少需要 $6 \leq 2^3$ ，subnet mask 的值 => 需要 Host ID 的前 3 個 bits 當作 subnet ID。所以新的 subnet mask 是由原本 Class B 的 default subnet mask 255.255.0.0 來改，改成 255.255.11100000.00000000 => 255.255.224.0 (2%)

subnet 的 ID 要從此 Class B Network ID 10100011.01101011. XXXXXXXX.XXXXXXXX 來改，需要 Host ID 的前 3 個 bits 當作 subnet ID。因此第 6 個 subnet ID (從 000 開始，第六個是 101) 為 10100011.01101011.10100000.00000000 => 163.107.160.0 (2%)

因此第一個可用 Host ID 為 10100011.01101011.10100000.00000001 = 163.107.160.1 (2%)

最後一個可用 Host ID 為 10100011.01101011.10111111.11111110 = 163.107.191.254 (2%)

->共有 $2^{13}-2=8190$ 個可用 Host ID (1%)

c. IP address, subnet mask, default gateway (3%)

4. Consider sending a 2480-byte datagram into a link that has an MTU of 820bytes, including 20-byte IP header. Suppose the original datagram is stamped with the identification number 1. List these segments in a table with their data lengths, IDs, flags and offsets. (表格中 data length, offset 每列 1 分。沒有解釋或不清楚，視狀況扣分，8%)

fragment	data lengths	ID	offset	flag
1				
.....				

Ans:

820 Bytes 內 IP data=820-20=800Bytes. 2480 byte-20 byte (datagram IP data)=2460Bytes, 分為 800, 800, 800, 60, 共 4 個 fragments, 加上 20bytes IP header 後，data length 為 820, 820, 820, 80.

fragment	data lengths	ID	offset	flag
1st	820 (1%)	1	0 (1%)	1
2nd	820 (1%)	1	800/8=100 (1%)	1
3rd	820 (1%)	1	100*2=200 (1%)	1
4th	80 (1%)	1	100*3=300 (1%)	0

(data length, offset, 每列 1 分。沒有解釋或不清楚，視狀況扣分，8%)

5. (a) Explain two network-layer functions (功能 1%, 說明 1%, 2% each, 4% total)
(b) What is the longest prefix matching? (2%)
(c) Explain four scheduling mechanisms used in router. (名稱 1%, 說明 1%, 2% each, 8% total)
(d) What four items can be returned from the DHCP server? (2% each, 8% total)
(e) What is the major motivation of NAT? (2%)
(f) Which field is used by NAT to support simultaneous connections with a single LAN-side address? (2%)
(g) How many addresses do IPv4 have? (2%) How many addresses do IPv6 have? (2%)
(h) Which network layer protocol does the traceroute program use? (2%) (32% total)

Ans:

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- (a) forwarding: move packets from router's input to appropriate router output (1%, 1%)
routing: determine route taken by packets from source to destination
 - (b) when looking for forwarding table entry for given destination address, use longest address prefix that matches destination address. (2%)
 - (c) FIFO (first in first out) scheduling: send in order of arrival to queue (1%, 1%)
Priority scheduling: send highest priority queued packet
Round Robin (RR) scheduling: cyclically scan class queues, sending one complete packet from each class (if available)
Weighted Fair Queuing (WFQ): each class gets weighted amount of service in each cycle (2%)
 - (d) allocated IP address on subnet (2%)
address of first-hop router for client
name and IP address of DNS sever
network mask
 - (e) range of addresses not needed from ISP: just one IP address for all devices (2%)
 - (f) 16-bit port-number in IP header (2%)
 - (g) 2^{32} IPv4 addresses; (2%)
 2^{128} IPv6 addresses (2%)
 - (h) ICMP (2%)
6. (a) Which command is used to release your current IP address? (2%)
(b) Which command is used to instruct your host to obtain a network configuration, including a new IP address? (2%)
(c) To see only the DHCP packets in Wireshark, which protocol is entered into the filter field? (2%)
(d) Which transport protocol are DHCP messages sent over? (2%)
(e) Before offering an IP address to a client, why should the DHCP server issues an ARP request for the offered IP? (2%) (10% total)

Ans:

- (a) ipconfig /release (2%)
- (b) ipconfig /renew (2%)
- (c) bootp (2%)
- (d) UDP (2%)
- (e) make sure the IP address is not already in use by another workstation. (2%)